

This Page Is Inserted by IFW Operations  
and is not a part of the Official Record

## **BEST AVAILABLE IMAGES**

Defective images within this document are accurate representations of the original documents submitted by the applicant.

Defects in the images may include (but are not limited to):

- BLACK BORDERS
- TEXT CUT OFF AT TOP, BOTTOM OR SIDES
- FADED TEXT
- ILLEGIBLE TEXT
- SKEWED/SLANTED IMAGES
- COLORED PHOTOS
- BLACK OR VERY BLACK AND WHITE DARK PHOTOS
- GRAY SCALE DOCUMENTS

IMAGES ARE BEST AVAILABLE COPY.

**As rescanning documents *will not* correct images,  
please do not report the images to the  
Image Problems Mailbox.**

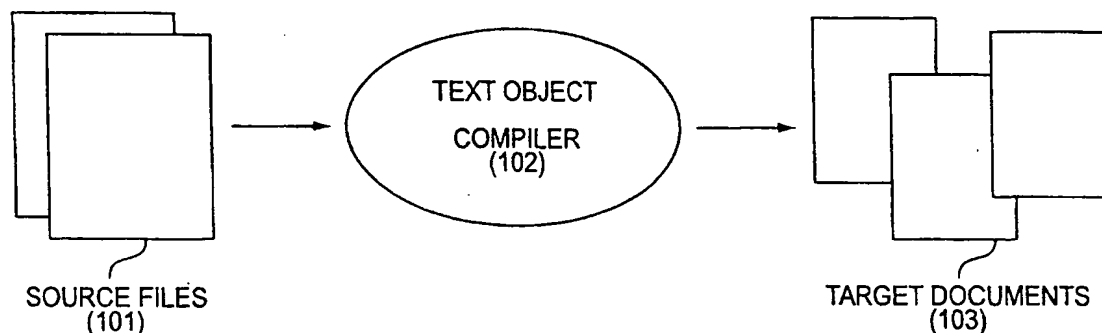
**THIS PAGE BLANK (USPTO)**



## INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(51) International Patent Classification <sup>7</sup> : <b>G06F 9/44</b>	<b>A1</b>	(11) International Publication Number: <b>WO 00/17748</b> (43) International Publication Date: 30 March 2000 (30.03.00)
<p>(21) International Application Number: PCT/US99/21940</p> <p>(22) International Filing Date: 22 September 1999 (22.09.99)</p> <p>(30) Priority Data: 09/159,304 23 September 1998 (23.09.98) US</p> <p>(63) Related by Continuation (CON) or Continuation-in-Part (CIP) to Earlier Application US 09/159,304 (CON) Filed on 23 September 1998 (23.09.98)</p> <p>(71) Applicant (for all designated States except US): NETCREATE SYSTEMS, INC. [US/US]; 2970 Fifth Avenue 320, San Diego, CA 92107 (US).</p> <p>(72) Inventor; and (75) Inventor/Applicant (for US only): DOMI, Dwayne, K. [US/US]; 2970 Fifth Avenue 320, San Diego, CA 92107 (US).</p> <p>(74) Agents: BURKE, John, E. et al.; Pillsbury Madison &amp; Sutro, LLP, 1100 New York Avenue, N.W., Washington, DC 20005 (US).</p>	<p>(81) Designated States: AE, AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CR, CU, CZ, DE, DK, DM, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, TZ, UA, UG, US, UZ, VN, YU, ZA, ZW, ARIPO patent (GH, GM, KE, LS, MW, SD, SL, SZ, TZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GW, ML, MR, NE, SN, TD, TG).</p> <p><b>Published</b> With international search report. Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.</p>	

(54) Title: TEXT OBJECT COMPILATION METHOD AND SYSTEM



## (57) Abstract

A Text Object Compiler and Language able to produce binary and text objects that are not machine language code. An object oriented computer language that produces target files of information in any text or binary format; files are defined by the programmer as "pages" and file locations are defined by the programmer as "targets". The compiler compiles the language to produce any variety of output, which include text formats (such as HTML, SGML, and other scripting languages) and binary formats (such as graphical pictures, binary data, or other multimedia information).

# FOR THE PURPOSES OF INFORMATION ONLY

Codes used to identify States party to the PCT on the front pages of pamphlets publishing international applications under the PCT.

AL	Albania	ES	Spain	LS	Lesotho	SI	Slovenia
AM	Armenia	FI	Finland	LT	Lithuania	SK	Slovakia
AT	Austria	FR	France	LU	Luxembourg	SN	Senegal
AU	Australia	GA	Gabon	LV	Latvia	SZ	Swaziland
AZ	Azerbaijan	GB	United Kingdom	MC	Monaco	TD	Chad
BA	Bosnia and Herzegovina	GE	Georgia	MD	Republic of Moldova	TG	Togo
BB	Barbados	GH	Ghana	MG	Madagascar	TJ	Tajikistan
BE	Belgium	GN	Guinea	MK	The former Yugoslav Republic of Macedonia	TM	Turkmenistan
BF	Burkina Faso	GR	Greece	ML	Mali	TR	Turkey
BG	Bulgaria	HU	Hungary	MN	Mongolia	TT	Trinidad and Tobago
BJ	Benin	IE	Ireland	MR	Mauritania	UA	Ukraine
BR	Brazil	IL	Israel	MW	Malawi	UG	Uganda
BY	Belarus	IS	Iceland	MX	Mexico	US	United States of America
CA	Canada	IT	Italy	NE	Niger	UZ	Uzbekistan
CF	Central African Republic	JP	Japan	NL	Netherlands	VN	Viet Nam
CG	Congo	KE	Kenya	NO	Norway	YU	Yugoslavia
CH	Switzerland	KG	Kyrgyzstan	NZ	New Zealand	ZW	Zimbabwe
CI	Côte d'Ivoire	KP	Democratic People's Republic of Korea	PL	Poland		
CM	Cameroon	KR	Republic of Korea	PT	Portugal		
CN	China	KZ	Kazakhstan	RO	Romania		
CU	Cuba	LC	Saint Lucia	RU	Russian Federation		
CZ	Czech Republic	LI	Liechtenstein	SD	Sudan		
DE	Germany	LK	Sri Lanka	SE	Sweden		
DK	Denmark	LR	Liberia	SG	Singapore		
EE	Estonia						

## TEXT OBJECT COMPILATION METHOD AND SYSTEM

### BACKGROUND OF THE INVENTION

#### Field of the Invention

5 This invention relates to the field of software compilers, and more particularly relates to a method of generating files of information from one or more source files.

#### Description of the Related Art

In March, 1989, the European Laboratory for Particle Physics or CERN (Conseil  
Europeen pour la Recherche Nucleaire) developed the World-Wide-Web (WWW, or  
10 simply, "the web"), an Internet-based computer network that allows users on one  
computer to access information stored on other computers through a world-wide network.  
With an intuitive user-interface, known as a web browser, the web rapidly became a  
popular way of transmitting and accessing text and binary information. Since then, there  
has been a massive expansion in the number of World-Wide-Web sites, and the amount of  
15 information placed on the web.

Information, in the form of electronic files, documents, images, sounds and other  
formats, forms the basis of internet and web content, and the key to creating a useful and  
meaningful web-site.

To place information on the web, the information must be stored in a binary or  
20 text format in a "file." Binary documents are saved in known formats that depend upon  
the information being stored. For example, two-dimensional pictures are often stored in  
"Joint Photographic Experts Group" (JPEG) or "Graphical Image Format" (GIF) standard  
formats. Audio files and moving images have other formats as well, such as "WAV,"  
"MOV," and "MPEG." For text documents, documents are stored in a HyperText Markup  
25 Language (HTML) format. The HTML format dictates the appearance and structure of a  
web text document, also referred to as a "web page."

Although these formats are required to create compatibility for web browsers,  
modifying web sites and updating information in these rigid formats is difficult and time  
consuming. For example, suppose every web page had a copyright notice on it. To  
30 update the copyright notice on every page, a web-site administrator would have to either  
change every page by hand, or use a method of global-search-and-replace. However,  
because of the non-uniform manner of some web-sites, a global-search-and-replace may

not work. More complicated web page changes, such as modifying small applications, known as "applets," are even more difficult. It would be much better if there was a single location or file that could be updated, and the change would be propagated to the entire web-site, or just the appropriate web pages. Very simply put, the problem of maintaining and generating large amounts of data, in any format, is difficult and highly time consuming.

Several solutions have been proposed, each has its problems.

Some web developers choose to generate web pages through a "what you see is what you get" (WYSIWYG) web-page editor. Such editors assemble web pages through a graphical interface, which makes designing pages simpler, but the results are limited because it does not solve the need to maintain the information. Using the above example, to update the copyright notice on every page, a web-site administrator would still have to edit the web pages individually, or the web-page editor program may use a method of global-search-and-replace.

Alternatively, simple pre-processor programs have been used to assemble HTML files. Such pre-processors allow web-page designers to pre-process documents and insert listed documents into a master document. For example, to include another listed document file called "foo.doc" into the master document, a web-page designer could type:

```
#include "foo.doc"
```

and the listed document would be included. While this allows fragments of common HTML code to be inserted into documents, as a web-site grows, and more pages are added to the site, the maintenance of such a system quickly becomes a logistical nightmare. Also, the fragments cannot be redefined at the point that they are included in a document. Moreover, such a system is limited strictly to text-based documents, and cannot handle binary forms of information.

U.S. Pat. No. 5,181,162, issued January 19, 1993 to Smith et al. entitled "Document management and production system," incorporated herein by reference, discloses a system of decomposing documents into logical components, which are stored as discrete "objects" in an object-oriented computational environment. The system relies on queries to a relational database which occur every time the document is printed, displayed electronically, or electronically transmitted. For a web site, which may transmit pages thousands of times per minute, this solution is a burden on the web server's computing resources. Consequently, the system would be slow, and of limited usefulness to such a high-demand environment. Similarly, the use of a relational database to deliver

pages of information on client machines has been attempted; while this provides dynamic construction of documents when they are delivered to the client machines, this solution also burdens the server's computing resources because page information would be constantly regenerated. Although caching generated pages may solve some of the computing resource problems, it creates a new problem because cached pages may be outdated.

Several related patents, U.S. Pat. No. 5,668,999, which issued September 16, 1997 to Gosling ("System and method for preverification of stack usage in bytecode program loops,"), U.S. Pat No. 5,692,047, issued to McManis ("System and method for executing verifiable programs with facility for using non-verifiable programs from trusted sources,"), and U.S. Pat No. 5,706,502, issued to Foley et al. ("Internet-enabled portfolio manager system and method,"), all incorporated herein by reference, also fail to solve the problem. Collectively, these patents disclose a method and system of verifying the integrity of computer programs written in a bytecode language to run applications remotely on a client workstation. While this solution may create dynamic client-machine applications, it does not solve the problem of maintaining information in a system.

What is needed is a more flexible way of handling both binary and text information that can produce files of different file formats and still be easy to maintain.

The invention, a Text Object Compiler method, allows users to abstract information, and produce information in virtually any file format.

Almost every contemporary computer is a register-based Von Neuman computer that responds to a machine language. These machine languages include instructions which operate on the contents of registers. Originally, computer software instructions were organized in terms of machine language operations. As computers became more complex, programming in machine language became difficult and increasingly cumbersome. Consequently, computer scientists abstracted machine instructions, creating higher-level languages, known as source languages, structured in terms of expressions and procedures. As software evolved, two strategies for converting source languages into instructions in machine language developed, interpreters and compilers.

An interpreter, written in the native machine language, configures the computer to execute programs written in one of the source languages. The primitive operators or commands of the source language are implemented as a library of subroutines written in the native machine language of the given machine. Interpreters read the source language,

one line at a time, and then perform the specified operation. A program to be interpreted, the source program, is represented as a data structure. The interpreter traverses this data structure, analyzing the source program. As it does so, it simulates the intended behavior of the source program by calling appropriate primitive operators from the library.

5           Instead of analyzing and translating the source program into machine language during execution, it is possible to perform these tasks before execution, enabling more efficient program execution. This alternate method of converting source languages into instructions is called compilation. The program that does the analysis of the source program and reduces the source program to machine language is called a compiler. As  
10 shown in FIG. 1, a conventional (i.e., prior art) compiler 2 for a given source language and machine translates computer source code 1 (i.e., a program written in a high level "computer language") into an object code 3, a program written in the computer's native language, referred to in the art as "machine language."

          Illustrated by FIG. 2, a conventional compiler 2 is composed of a lexical analyzer  
15 10, a parser 20, and code generator 30. A lexical analyzer 10 takes computer source code 1 and divides the code into lexical tokens. Such lexical tokens can be based on instructions or other keywords in the relevant high level computer language. A parser 20 takes the tokens and groups them together logically based on the relationships established by the source language and the computer source code 1. Lastly, a code generator 30 takes  
20 the relationships established by the parser 20 and translates them into an executable computer object code 3 in computer machine language.

          Conventional compilers are well known in the prior art, such as U.S. Pat. No. 5,560,015 ("Compiler and method of compilation" issued to Onodera on September 24, 1996), U.S. Pat. No. 5,442,792 ("Expert system compilation method" issued to Chun on  
25 August 15, 1995), and U.S. Pat. No. 5,768,592 ("Method and apparatus for managing profile data" issued to Chang on June 16, 1998) all incorporated herein by reference.

          Conventional interpreters and compilers convert high-level computer source code into object code to be executed on a computer. In effect, the interpreter and the compiler allow computer programmers to write computer programs at a higher level of abstraction,  
30 and generate object code.

#### SUMMARY OF THE INVENTION

          The invention, a Text Object Compiler (TOC) method and system, applies this same level of abstraction to information as a conventional compiler applies to computer



programs. A user designs abstract source code which is compiled into a file or a plurality of files, which is not object machine language code. Instead, the TOC produces information in virtually any information format, as text or binary files. The TOC reads one or more source files written as a Text Object Language (TOL) in ASCII text and  
5 processes the source files into one or more output files in any document format. The compilation process reorganizes the information in the source files into output document formats, and may contain compile-time utility commands to facilitate the document generation process.

In the first embodiment, a lexical analyzer tokenizes a source input written in TOL  
10 regular expressions to produce a token representing the source input. Such TOL expressions may contain variables, functions, and classes. A parser determines the relationships between the tokens so that a page generator can evaluate the tokens and the relationships between the tokens to generate an output. The resulting output may be written as a file at a target location specified by the source input.

15 In another aspect of the invention, the source input is lexically analyzed to produce tokens representing regular expressions of the source input. The regular expressions are written in the Text Object Language and may include variables, functions, and classes. The regular expressions are parsed to determine the relationship between the tokens. For instance, any class relationship between tokens is determined within this step.  
20 Finally the tokens and relationships are evaluated to generate a non-executable output. The non-executable output may then be written to a file. The file location may be specified by the original source input as a specific target location.

In the preferred embodiment of the present invention, the source input, written regular expressions of a Text Object Language, is initially read. The regular expressions  
25 contain page definitions used to determine the output of the process, and may additionally contain variables, functions, target locations and object-oriented classes. Once read, the source input is lexically analyzed to produce token representations of the regular expressions, which includes tokens generated from the page definitions. The tokens are parsed to determine the relationship between the tokens, and the resulting relationship is  
30 constructed in computer memory. Each variable and function is evaluated and their value is determined. The determined value replace their corresponding variable or function token in computer memory. The non-executable file based on the computer memory representation of the page tokens is then written to a file at a target location specified

within the source input.

### BRIEF DESCRIPTION OF THE DRAWINGS

The foregoing and other objects, features, and advantages of the present invention will be better understood in view of the following detailed description made in conjunction with the accompanying drawing in which:

FIG. 1 diagrams an overview of the conventional compiler process (prior art);

FIG. 2 illustrates the basic components of a conventional compiler (prior art);

FIG. 3 diagrams an overview of the Text Object Compiler (TOC) process;

FIG. 4 illustrates the basic components of a Text Object Compiler;

FIG. 5 is a flowchart of the lexical analysis and parsing subprocesses used by the Text Object Compiler.

FIG. 6 is an inheritance diagram of the classes of the Text Object Language source file listed in Table 3;

FIG. 7 is an inheritance diagram showing the relationship of class variables and the classes;

FIG. 8 is an inheritance diagram showing the relationship of class functions and the classes;

FIG. 9 is an inheritance diagram consolidating the classes with their functions and variables;

FIG. 10 is an inheritance diagram showing the relationship of the defined pages and the classes;

FIG. 11 diagrams the baseclass information used to generate the target document "first.txt";

FIG. 12 diagrams the baseclass, and myclass information used to generate the target document "second.txt";

FIG. 13 diagrams the baseclass, myclass, and newclass information used to generate the target document "third.txt"; and

FIG. 14 is a flowchart detailing the page generation subprocess.

### DETAILED DESCRIPTION OF THE INVENTION

The Text Object Compiler uses a variety of existing programming and compiler methods that are commonly used in programming languages to create executable files. The unique aspect of TOC is that it is applied to creating non-executable files, which are

referred to as "target documents." As shown in FIG. 3, source files 101, written in a Text Object Language (TOL) are compiled by the TOC 102 to produce target documents 103 as output. The target document locations are definable by the programmer, and the location is referred to as a "target location." Target documents 103 can be any type of format, and can even produce other source files used by other compilers or programs. The TOC 102 actually knows nothing about the format of the target documents 103; the target document format is solely up to the programmer.

Reference will now be made in detail to the preferred embodiments of the invention, examples of which are illustrated in the accompanying drawings. While the invention will be described in conjunction with the preferred embodiments, it will be understood that they are not intended to limit the invention to those embodiments. On the contrary, the invention is intended to cover alternatives, modifications and equivalents, which may be included within the spirit and scope of the invention as defined by the appended claims.

#### The Text Object Language (TOL)

In the preferred embodiment, the source file or files are written in a Text Object Language, which when compiled or interpreted will result in at least one target file. A listing of some of the TOL operators is provided in Table 1.

<u>Operator</u>	<u>Description</u>
<code>:=</code>	The equality operator tells the compiler to define the keyword on the left-hand-side of the operator as the value on the right-hand-side of the operator. For example, the usage " <code>length := 5</code> " would define the variable " <code>length</code> " with the value of 5.
<code>//</code>	This operator indicates to the compiler the presence of a comment-line. The compiler ignores the remainder of the line.
<code>#include := filename</code>	Imports a source file named " <code>filename</code> " for compilation.
<code>class classname := baseclass</code>	Defines a page class. <i>Classname</i> must be specified. <i>Baseclass</i> , if omitted, is assumed to be the base class for TOL. A class inherits all of the variables and functions of its base class. Classes can be <i>public</i> , <i>private</i> , or <i>protected</i> . A class can declare other classes as <i>friends</i> ; a friend class is given public access to all of the declaring class's variables and functions. Within a class, variables or functions can be declared friends.

**Cleartargets := directory****Clears all targets previously inline.****SUBSTITUTE SHEET (RULE 26)**

<code>Debug := {on/off}</code>	Enables/Disables debug while compiling. The default is to stop the debugger from source code.
<code>func(classname) := name (var1, ... , varN = value) endfunc</code>	<p>Begins a new function. If <i>classname</i> is omitted, then the class is assumed to be the base class. The function name, <i>name</i>, must be unique for a class. The function can contain zero or more variables, <i>var1</i>, <i>var2</i>, etc.</p> <p>Functions can be overloaded, with several functions of the same name, but with different number of variables; overloaded functions must contain a unique number of variables.</p> <p>Functions can be <i>public</i>, <i>private</i>, or <i>protected</i>.</p> <p>Functions can be made virtual, forcing them to be defined in a derived class. Once a function is declared virtual, all derived class instances of the function are virtual.</p> <p>Functions can contain a default value, for example, <i>var1=default</i>.</p> <p><i>Endfunc</i> ends a function section.</p>
<code>lfor := {on/off}</code>	Enables/Disables output file wrapping.
<code>page(classname) := filename endpage</code>	<p>Begins a new page section. If <i>classname</i> is omitted, then the class is assumed to be the base class. The compiler builds an output page for each page/endpage section. Within this section, all variables and functions are resolved.</p> <p><i>Filename</i> is the fully qualified location of the output file.</p> <p><i>Endpage</i> ends a page section.</p>
<code>targets (classname) endtargets</code>	<p>Targets allow pages output be directed to different or multiple locations. Example targets include any media or memory storage device, such as disk drives, networked drives, FTP locations, memory cards.</p> <p><i>Endtargets</i> ends a target section.</p>
<code>vars (classname) endvars</code>	<p>Begins a new variables section. If <i>classname</i> is omitted, then the class is assumed to be the base class.</p> <p>Variables can be <i>public</i>, <i>private</i>, or <i>protected</i>.</p> <p>Variables can be made virtual, forcing them to be defined in a derived class. Once a variable is declared virtual, all derived class instances of the variable are virtual.</p>

Table 1. TOL Operators

The TOL is similar to other programming languages, in that it has variables, classes, functions and subroutines, but uses a language syntax recognized only by the TOC 102.

In addition, to the TOL operators, a number of compile-time utility commands exist to facilitate the document generation process

<u>Utility Command</u>	<u>Description</u>
<b>beep</b> := <i>length</i>	Causes the system to sound beep with duration of <i>length</i> in milliseconds. The default duration is 100 milliseconds.
<b>chdir</b> := <i>path</i>	Changes the current directory to <i>path</i> . Any reference to a path or filename that does not include a drive letter will default to the current drive. "chdir" is executed by the compiler inline.
<b>Chdrive</b> := <i>driveletter</i>	Changes the current drive. Any reference to a path or filename that does not include a drive letter will assume the current drive. "chdrive" is executed by the compiler inline.
<b>copy</b> := <i>source, destination</i>	Copies source to destination when the compiler reaches this inline command.
<b>exec</b> := <i>program</i>	Runs the specified application in program.
<b>kill</b> := <i>filespec</i>	Deletes the files specified in <i>filespec</i> .
<b>md</b> := <i>path</i>	Creates the directory <i>path</i> .
<b>rd</b> := <i>path</i>	Removes the directory <i>path</i> .

Table 2. TOL Compile-Time Utility Commands

An exemplary source file 101 written in the Text Object Language of Table 1 can be seen in Table 3.

```
//define the class structure and relationships
class myclass := baseclass
class newclass := myclass

//variable definitions

//baseclass variables
vars
    title := This is my document title
endvars

//variables for myclass pages only
vars(myclass)
    title := This is the document title for myclass
endvars

//functions

//baseclass function example
func:=myfunc (var1, var2)
    <H1>var1</H1>
    <H2>var2</H2>
endfunc

//function for newclass pages only
func(newclass):=myfunc(var1, var2)
    <center>
        var1<br>
        var2
    </center>
endfunc

//target documents
page:=first.txt
    myfunc(title, This is a base class example)
endpage

page(myclass):=second.txt
    myfunc(title, This is a myclass example)
endfunc

page(newclass):=third.txt
    myfunc(title, This is a newclass example)
endpage

//target locations
targets
    Local Drive := c:\web
    LAN Drive := n:\web
    Live Site := ftp://www.netcreate.com/web/html
Endtargets
```

**Table 3. Example Source File written in the Text Object Language ("mysource.txt")**



As illustrated in the example Text Object code source file 101 of Table 3, there are five primary components of TOL: classes, variables, functions, pages, and targets.

Although the concepts of classes, variables, and functions exist in prior art computer languages, the additional concepts of pages and targets exist in the Text Object Language.

5 TOL classes are similar to and share many of the elements of common Object Oriented Programming (OOP) classes. Classes allow a document programmer to organize document sections into objects that can be reused throughout the source files and applies to any of the target files. Specifying classes is optional, since a default class, or "base class," is always assumed. A "derived" or "child" class inherits all of the  
10 variables and functions of its base or "parent" class.

Classes can be *public*, *private*, or *protected*. Public classes allow their functions and variables to be redefined by other classes. By default, all classes are public. Private classes allow their functions and variables to be redefined only by other member or friend classes. Protected classes allow its variables and function to be used only by member  
15 functions, friends of the class in which it is declared, and by member functions and friends of classes derived from the protected class. In addition, a class can declare other classes as *friends*; a friend class is given public access to all of the declaring class' variables and functions.

Variables allow the source file programmer to represent elements of a target  
20 document by reference, and use the reference to create sections or target documents rather than using the actual data. Like classes, in the preferred embodiment, variables can be can be *public*, *private*, or *protected*. Variables can also be made virtual, forcing them to be defined in a derived class; however, once a variable is declared virtual, all inherited class instances of the variable are virtual. Note that no virtual variables of the class may  
25 exist within the program until the virtual variable is defined by the derived (child) class.

A function is a convenient way to encapsulate some computation, which can then be used without worrying about its implementation. Functions allow programmers a conceptual way to abstract a recurring procedure without worrying about the details. Functions are similar to typical programming subroutines. Like classes and variables,  
30 functions can be *public*, *private*, or *protected*.

Function name overloading allows multiple function instances that provide a common operation on different argument types to share a common name. Functions can be overloaded, with several functions sharing the same name, but each having a different

number of variables. Each overloaded function must have a unique number of variables, which allows the compiler to distinguish between each instance of the overloaded function. Functions can be made virtual, forcing them to be defined in a derived (child) class. Once a function is declared virtual, all derived class instances of the function are virtual. A virtual baseclass function is also virtual in the derived class if inherited by the derived class; such a function is treated as an abstract class, and no objects of the class may exist within the program until the function is defined by a derived class.

Pages are unique to the Text Object Language; page parameters instruct the Text Object Compiler 102 how to combine or parse source files 101 into the actual individual target documents 103. A page defines the starting and ending point of a resulting target document 103, and the contents of the target document 103.

Targets are also unique to the Text Object Language. Target parameters define the target location; once defined, the target parameters instruct a Text Object Compiler 102 on where to place the target documents. This location is referred to as the "target location." The target location may be local to the computer running the TOC 102, or at a remote location that can be accessed over a computer network by the computer. If the source files 101 define multiple target locations with the target parameter, the TOC 102 will produce identical target documents 103 at each target location. Multiple targets are useful for creating experimental output, creating backup files for redundancy purposes, and updating main/production server files. For example, a programmer may define two targets to create a primary web-site and its "mirror" web-site at an alternate location. If target parameters are omitted from the Text Object code 101, the target documents 103 will be created in a default local location.

#### The Text Object Compiler (TOC)

The Text Object Compiler 102 performs the compilation of the text object language source files 101, resulting in target documents 103 defined by the pages parameter as output at a target location defined by the targets parameter. As discussed, target documents 103 may be in any format; note that this distinguishes the TOC 102 from prior art software compilers that which only produce object machine language code, i.e. executable files. Note, however, that programmers define the output format of the files with their source program code 101.

Attention will now be given to the TOC structure and method.

The TOC 102 is similar in structure to a conventional compiler. Like a conventional

conventional compiler, the TOC contains a lexical analyzer 10 and a parser 20. However unlike the TOC, a conventional compiler, shown in FIG. 2, feeds parser output into a computer code generator 3, to generate executable computer object code 3. As illustrated in FIG. 4, in a Text Object Compiler 102, the parser 20 output is presented to a page generator 200 to produce the target documents 103 as output.

The TOC lexical analyzer 10 examines expressions in a similar fashion to a conventional compiler lexical analyzer. This division into units, known as "tokens," is a process known in the art as "lexical analysis." Essentially, the lexical analyzer looks for regular expressions. A regular expression is a pattern description using the computer language. The lexical analyzer performs as many regular expression matches as possible, and attempts to classify the text of the entire source file into tokens. In the Text Object Language, the expressions may include variable names, function names, class names, target locations, page definitions, constants, strings, operators, punctuation, and so forth. For example, when compiling the source file in Table 3, the compiler initially classifies each instance of a known operator (as listed in Table 1) as a known token. However, if the word or expression is unknown to the compiler, it too is still tokenized, but its value or relationship must still be determined by the parser.

As the input is divided into tokens, the compiler must establish the relationship between the tokens. The Text Object Compiler needs to find the expressions, statements, declarations, blocks, functions/procedures, class structures, and pages in the program, a process known as "parsing." The list of rules that define the relationships that the compiler understands is called grammar. The grammar of an exemplary Text Object Language is shown above in Table 1.

The Text Object Compilation process is best explained by example. An existing source file 101, such as the example in Table 3 is written in the Text Object Language. The compiler reads the source file, as illustrated in step 250 of FIG 5. In its process of compiling the source code, a compiler performs two tasks over and over: a.) dividing the input source code into meaningful units (step 260), and b.) discovering the relationship between the units (step 270). These two processes are respectively called "lexical analysis" (step 260) and "parsing" (step 270). If the parser cannot determine the relationship of the token, it next determines whether the end of the source file has been reached, step 280. If the end of the source file has been reached (step 280), the undetermined tokens are an error in either syntax or usage, and an error is reported, step

282. If the end of the source file has not been read, the compiler loops back to step 250, and reads the source file. Similarly, if the parsing of step 270 is successful, and the entire file has not been read, as determined by step 284, the compiler continues to read more lines of the source file, step 250.

5 An example of the lexical analysis and parsing steps are as follows. The compiler initially reads the first line of Table 3, step 250. Each word is tokenized, and matched against a known set of regular expressions, such as the TOL Operators. The first known operator, the comment operator ("//") is identified, step 260. As defined by the implementation of this TOL grammar, the remainder of the line is determined to be a  
10 comment, and the compiler ignores the remainder of the line, step 270. Since the end of the source file has not been reached, as determined by step 284, the compilation process continues, and the compiler reads the next line of the source file, step 250.

The lexical analyzer notes the presence of four tokens on the second line, the words "class," "myclass" "!=" and "baseclass." Step 260. Two of these tokens,  
15 "class" and the equality operator ("!=") are identified as operators, and a third token, "baseclass" is identified as the "baseclass" keyword, which defines the TOL base class. The token "myclass" is initially unknown by the lexical analyzer. The token information is forwarded to the parser, which realizes that the source file defines a child class "myclass" which descends from the TOL baseclass, step 270. The parser constructs a  
20 memory table, memory tree, or equivalent memory structure to categorize the class structure. The process is repeated with the next line, resulting in a class inheritance relationship depicted by FIG. 6. Class myclass 310 is derived from the base class 300, and class newclass 320 is a "child" class derived from the "parent" class myclass 310. The memory tree is expanded to reflect the newclass class.

25 The compiler processes the next several lines of Table 3, which consist of variable definitions for the variable "title." The variables are parsed and stored in the memory tree, linked to their appropriate class definition, as shown by FIG 7. The baseclass 300 is associated with a variable "title" 301. Similarly, myclass 310 is associated with a different definition for another variable called "title" 311.

30 FIG 8 illustrates the relationships of the functions declared with their defined classes. The code in Table 3 defines a "myfunc" function that is different for the baseclass 300 and newclass 320; consequently, a myfunc 302 is associated with the baseclass 300, and a different myfunc 322 function is associated with newclass 320.

FIG 9. consolidates the inheritance diagrams with their related variables and functions. Baseclass 300 has both a variable, title 301, and a function, myfunc 302. The class myclass 310 also has a variable, title 311, and since it does not have a definition for myfunc, it inherits the function definition for myfunc 312 from the baseclass 300 definition of myfunc 302. Similarly, the class newclass 320 does not have a value for the "title" variable, and thus inherits its definition for title 321 from the myclass title definition 311. Newclass 320 does have its own definition for the function myfunc 322, and this is also reflected in the inheritance diagram.

The lexical analysis and parsing process is repeated for both the target document and target location sections of the code. As shown in FIG. 10, the target document "first.txt" 400 is of the baseclass 300, "second.txt" 410 is of class myclass 310, and "third.txt" 420 is of class newclass 320. Each of the three target documents consist of a single function call to the appropriate class function "myfunc."

Once all the source files have been tokenized and parsed to known values, the established token and relationship information is passed to the compiler page generator, step 286.

As shown previously in FIG. 4, the compiler page generator 200 creates each target document based on the relationships and tokens forwarded from the parser 20, replacing variables with their appropriate values, evaluating function calls, and substituting the resulting information into the page table shown in FIG. 10.

The page generator sub-process is elaborated in FIG. 14. The token relationship information is passed to the compiler page generator, step 286. For each page class, the variables are replaced with their respective definitions, step 288. In a simple embodiment, this can merely be the substitution of the value into each memory table location where the variable appears. Each function call for every page class is then evaluated, step 290. The existence of each named target location is verified; if the location, such as a directory, does not exist, it may be created at this time by the compiler, step 292. Each page, corresponding to a target document, is then written at each target location, step 294. Lastly, the write is verified by the compiler, step 296.

For example, as illustrated in FIG. 11, the output for "first.txt" 400 is generated by noting the appropriate class, baseclass 300, which defines the functions and variables used in generating the page. Table 3 defines "first.txt" as a page generated by a function call to "myfunc" using the "title" variable and "This is a base class example" as the input.

Since "first.txt" is of class baseclass, the definitions for "title" and "myfunc" are taken directly from the baseclass. The results are shown in Table 4.

```
<H1>This is my document title</H1>
<H2>This is a base class example</H2>
```

**Table 4. Compiled output for "first.txt"**

FIG. 12 continues the compilation for "second.txt" 410, which is of class "myclass" 310. The definitions for "title" is taken directly from class myclass 310. The definitions for "myfunc" would normally also be taken from class myclass 310. However, since "myfunc" is not defined for myclass 310, the myfunc function definition for myclass' parent class, baseclass 300, is used. The compiled results for "second.txt" 410 are shown in Table 5.

```
<H1>This is the document title for myclass</H1>
<H2>This is a myclass example</H2>
```

**Table 5. Compiled output for "second.txt"**

FIG. 13 continues the compilation for "third.txt" 420, which is of class "newclass" 320. The definitions for "title" and "myfunc" are normally taken directly from class newclass 310. However, since the "title" variable is not defined for newclass 320, the "title" variable definition for newclass' parent class, myclass 310, is used. Since "myfunc" is defined for the class newclass 310, the newclass "myfunc" definition is used. The compiled results for "third.txt" 420 are shown in Table 6.

```
<center>
  This is the document title for myclass<br>
  This is a newclass example
</center>
```

**Table 6. Compiled output for "third.txt"**

Once the compiler generates each page in memory, each page is written as a target document at each target location. The compiler may optionally create previously non-existing target locations, and verify the writing at the target locations; in its most preferred embodiment, the TOC performs both actions, reporting a warning message if a target location is not created, or an error message if a problem in writing the target document occurs.

CLAIMS

We Claim:

1. A computer comprising:
  - 5 a. a lexical analyzer that tokenizes a source input to produce tokens representing regular expressions of the source input;
  - b. a parser that determines relationships between the tokens;
  - c. a page generator that evaluates the tokens and the relationships between the tokens to generate an output, the output being non-executable.
- 10 2. A computer of claim 1 wherein the page generator writes the output as a file.
3. A computer of claim 2 wherein the page generator writes the file at a target location specified within the source input.
4. A computer of claim 1 wherein the regular expressions include variables.
5. A computer of claim 1 wherein the regular expressions include functions.
- 15 6. A computer of claim 1 wherein the regular expressions include classes.
7. A computer of claim 1 wherein the regular expressions include variables, functions, and classes.
8. A method of operating a computer system to compile a Text Object Language comprising:
  - 20 lexically analyzing a source input to produce tokens representing regular expressions of the source input;
  - parsing the tokens to determine relationships between the tokens;
  - evaluating the tokens and the relationships between the tokens to generate a non-executable output.
- 25 9. A method of claim 8 further comprising:
  - writing the non-executable output to a file.
10. A method of claim 9 wherein the file is written at a target location specified within the source file.
11. A method of claim 8 wherein the regular expressions include variables.
- 30 12. A method of claim 8 wherein the regular expressions include functions.
13. A method of claim 8 wherein the regular expressions include classes.
14. A method of claim 8 wherein the regular expressions include variables, functions, and classes.

15. A method of operating a computer system to compile a Text Object Language comprising:

reading a source input containing regular expressions of the Text Object Language, wherein the regular expressions of the text object language include variables,  
5 functions, and page definitions;

lexically analyzing the source input to produce tokens of the regular expressions, wherein the tokens include page tokens;

parsing the tokens to determine relationships between the tokens;

constructing a representation of the tokens and their relationships in computer  
10 memory;

evaluating the tokens that represent variables and functions to determine their evaluated values;

replacing the computer memory representation of the variable tokens and the function tokens with their evaluated values;

15 writing non-executable files based on the computer memory representation of the page tokens.

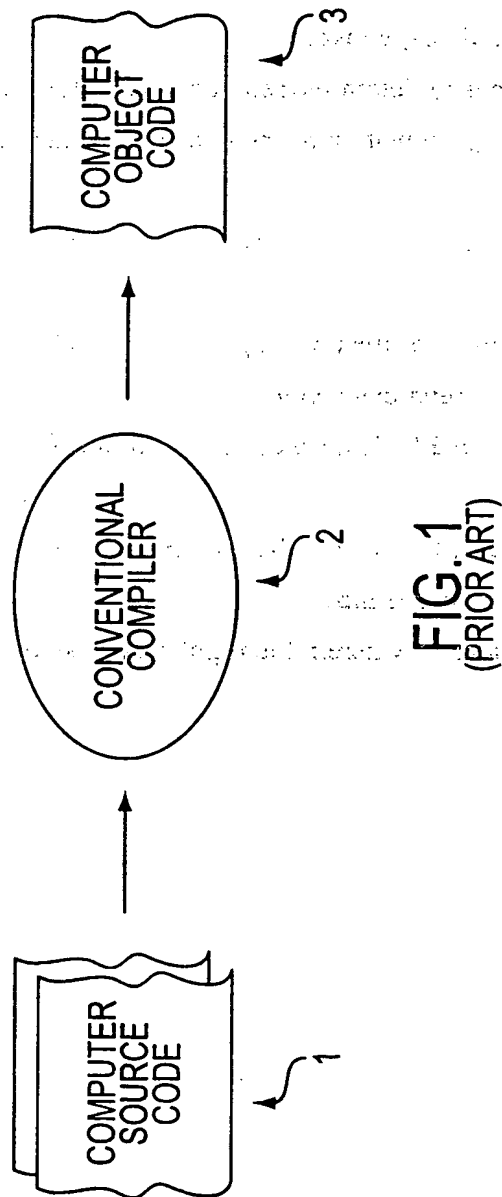
16. A method of claim 15 wherein the non-executable files are written in a target location specified by the source input.

17. A method of claim 15 wherein the regular expressions include classes.

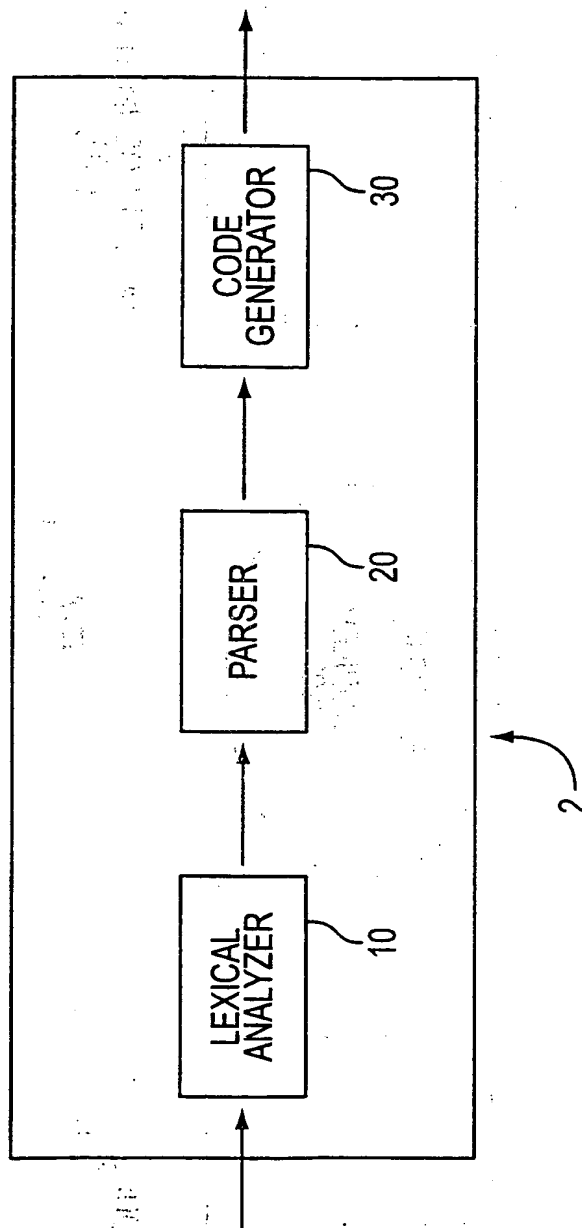
20



1/14



2/14

**FIG. 2**  
(PRIOR ART)

3/14

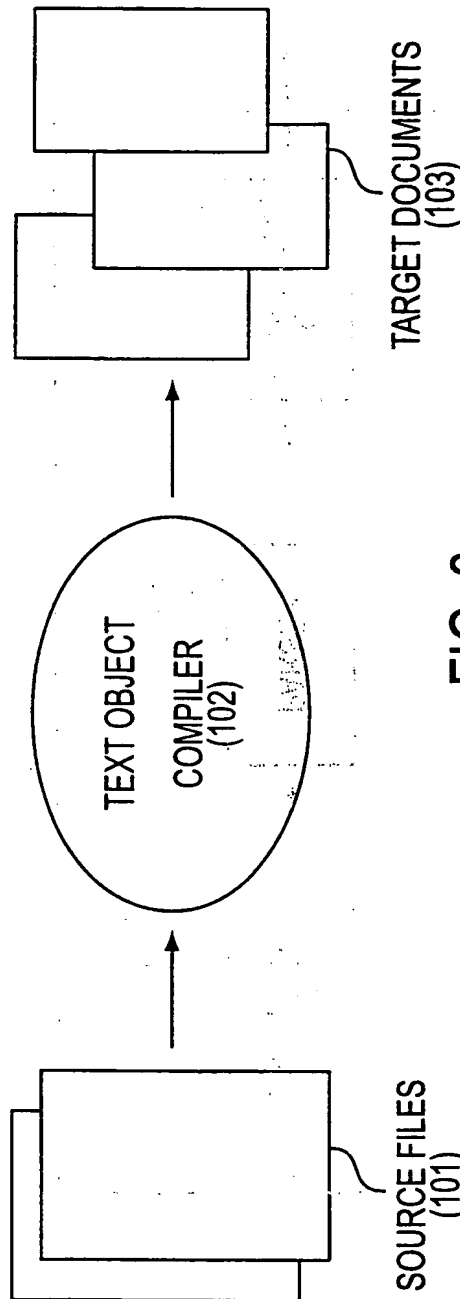


FIG. 3

4/14

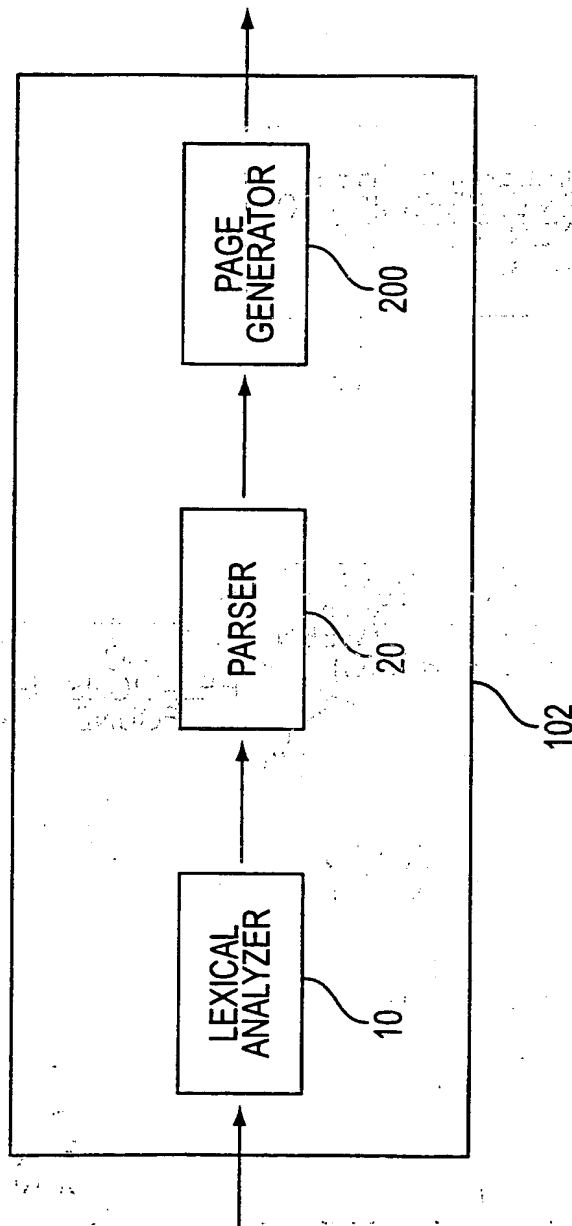


FIG. 4

5/14

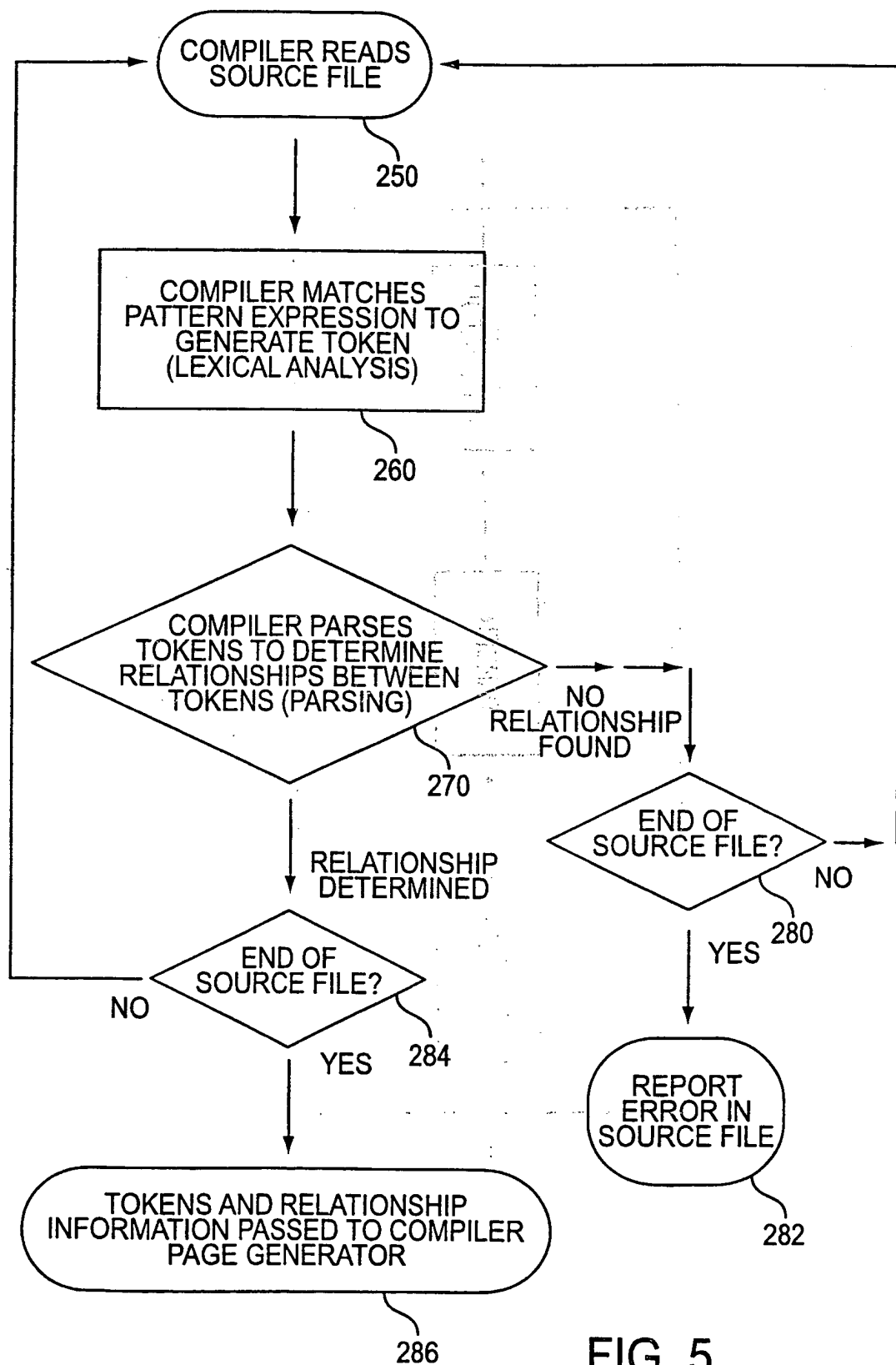


FIG. 5

6/14

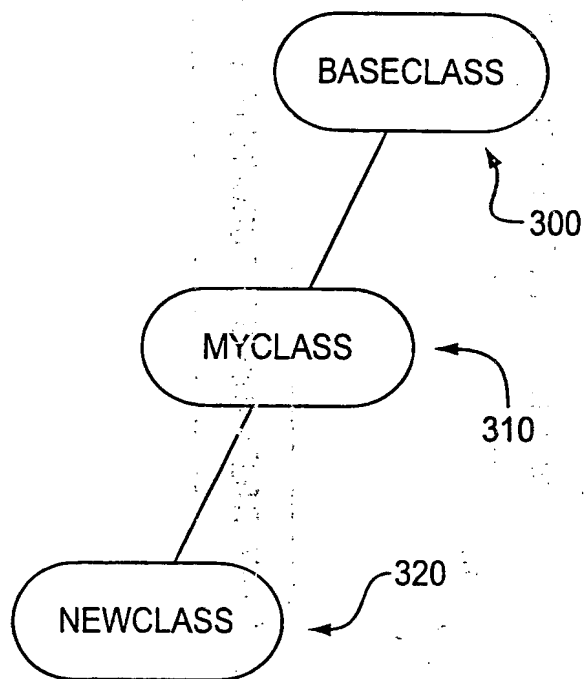


FIG. 6

7/14

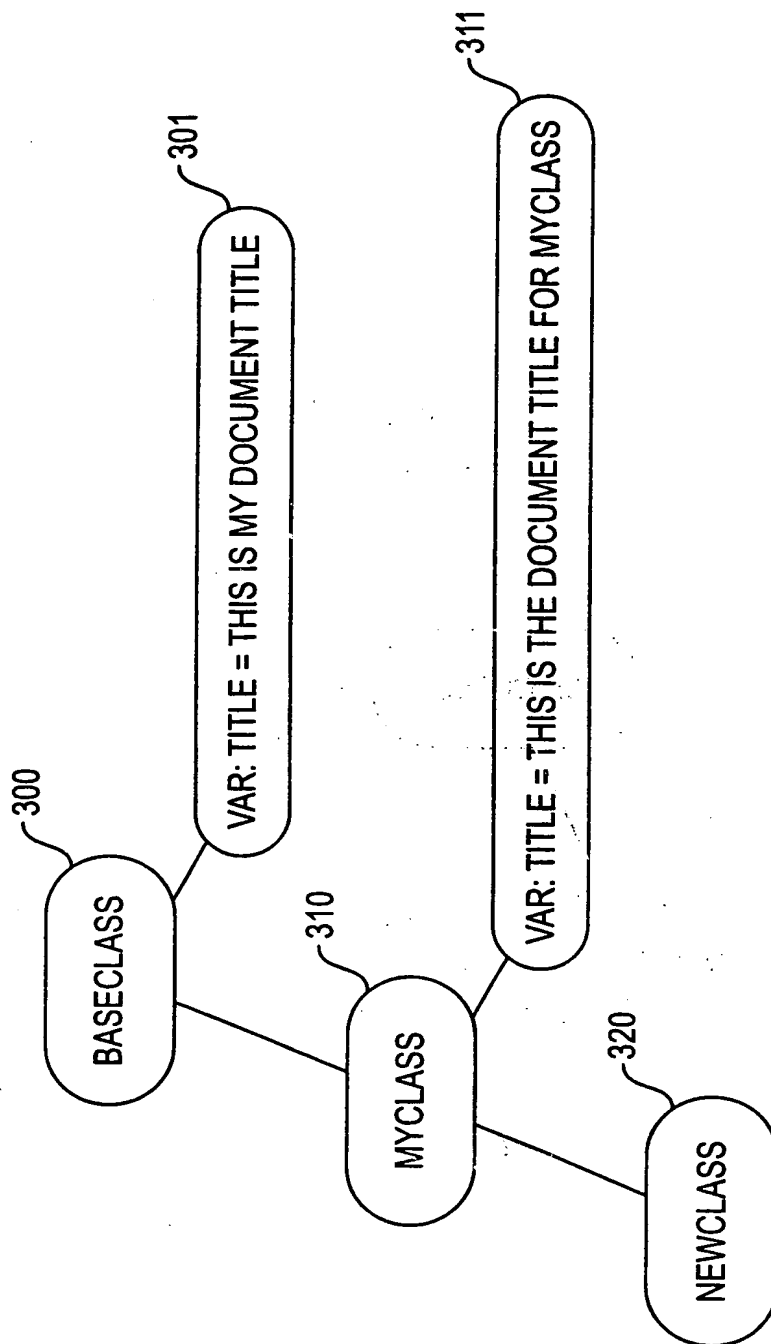


FIG. 7

8/14

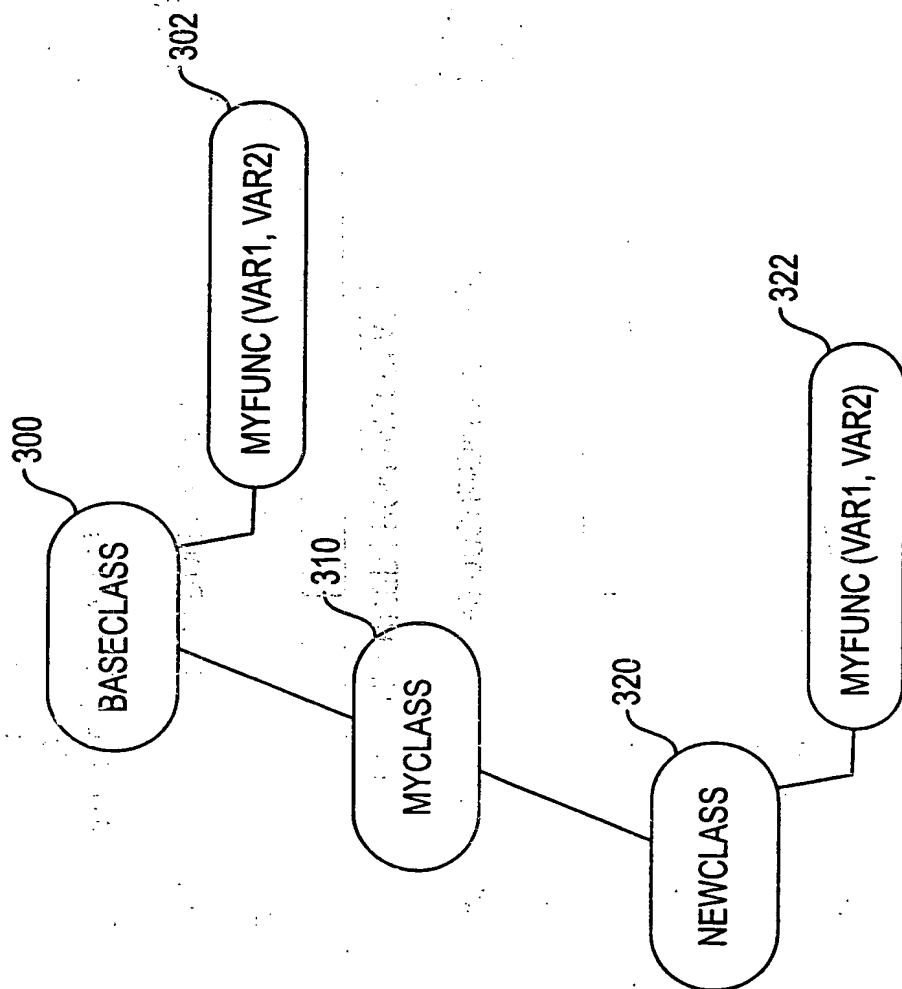


FIG. 8



9/14

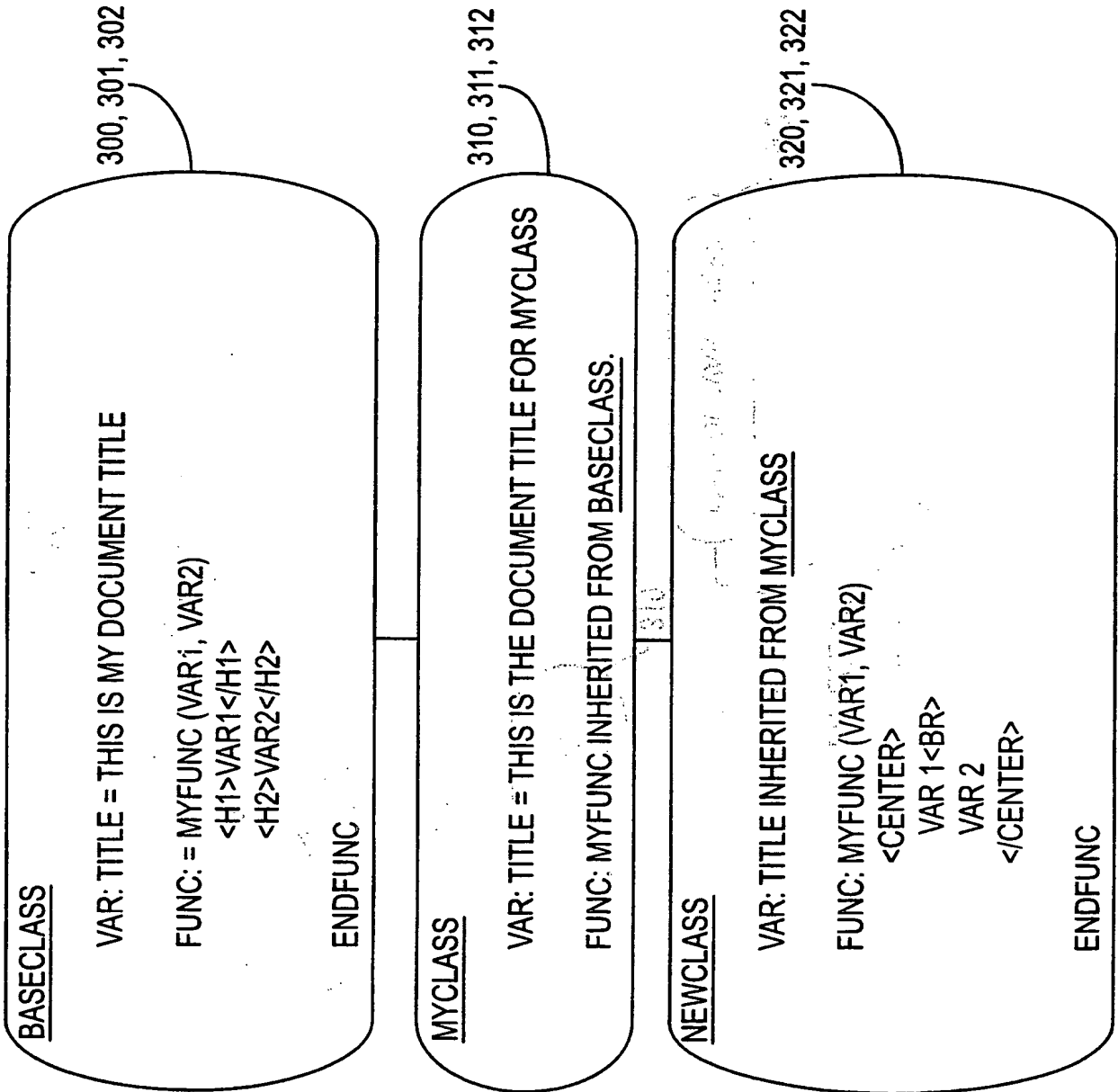


FIG. 9

10/14

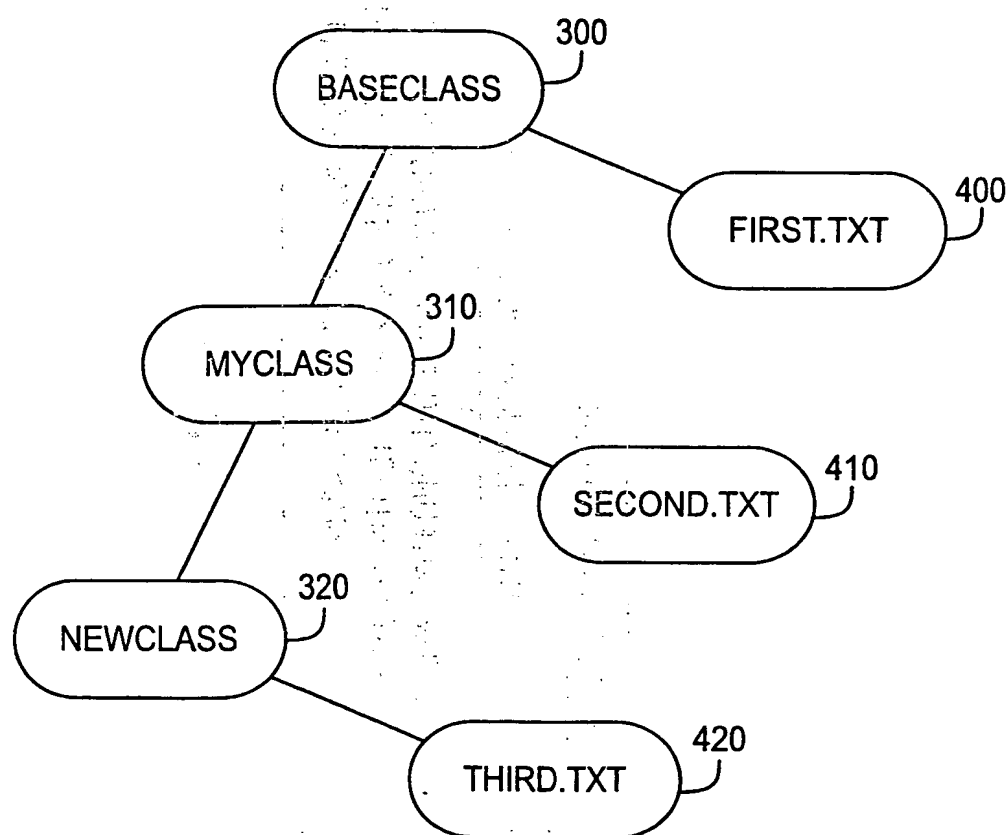


FIG. 10

11/14

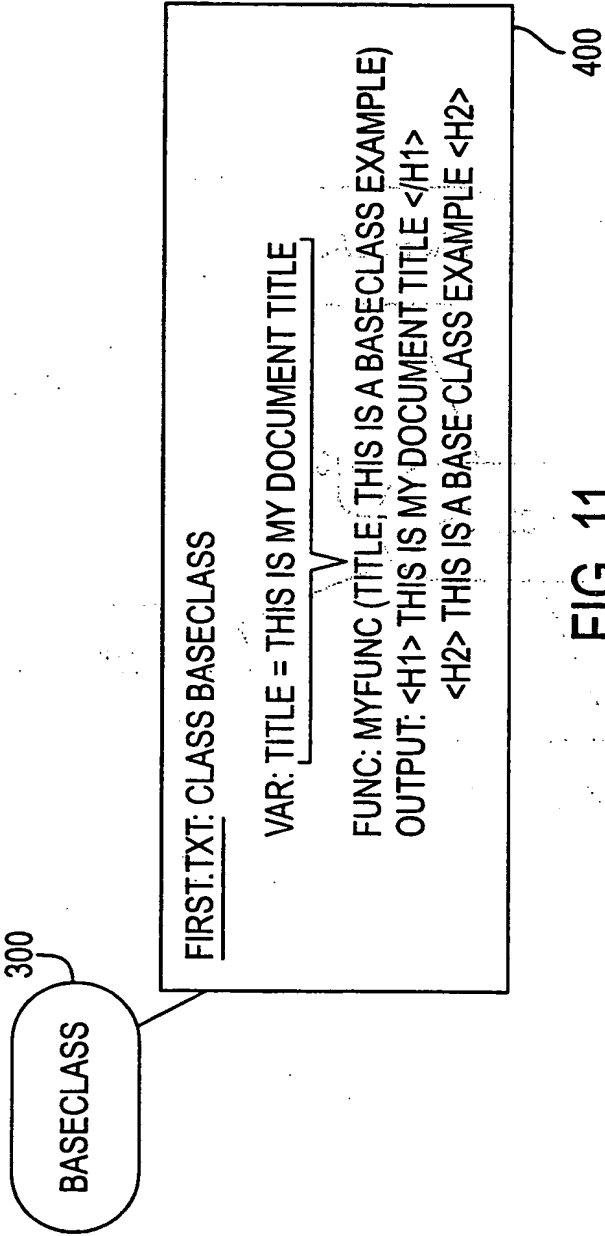


FIG. 11

12/14

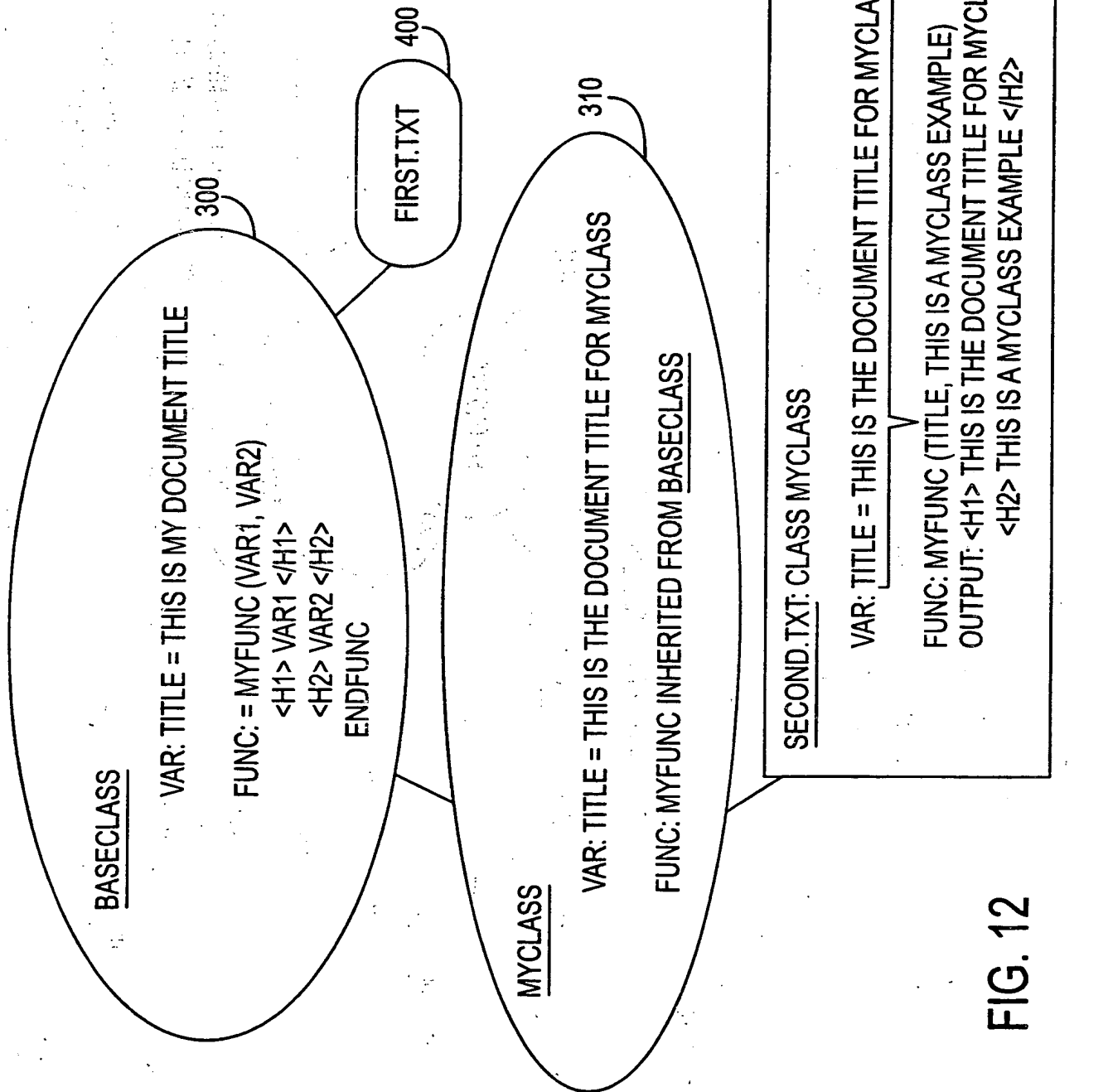


FIG. 12

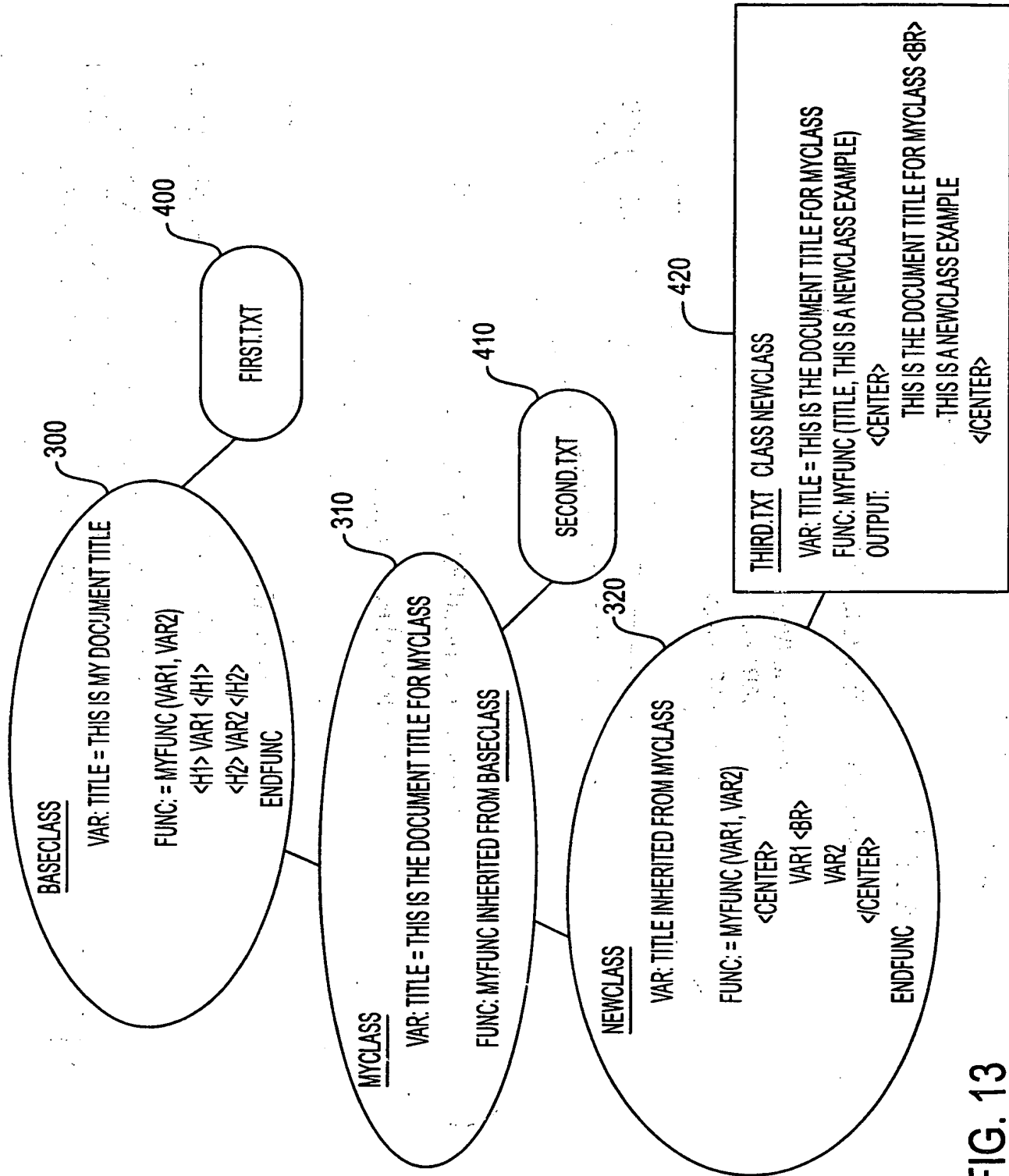


FIG. 13

14/14

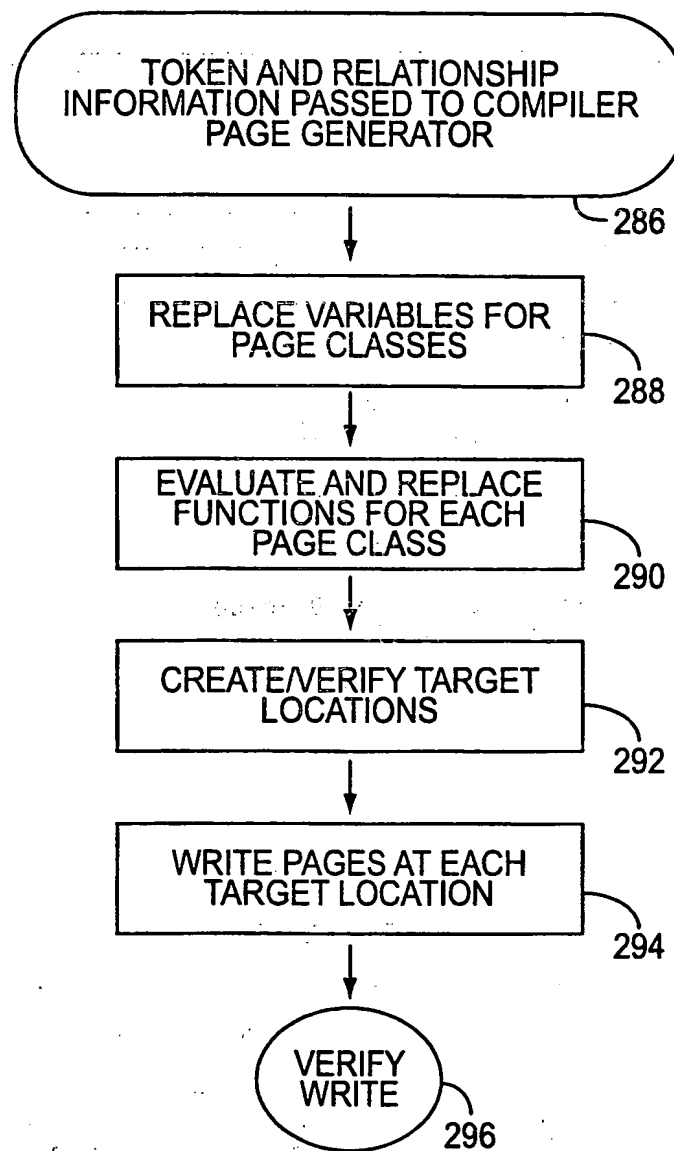


FIG. 14

# INTERNATIONAL SEARCH REPORT

International Application No

PCT/US 99/21940

## A. CLASSIFICATION OF SUBJECT MATTER

IPC 7 G06F9/44

According to International Patent Classification (IPC) or to both national classification and IPC

## B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 7 G06F

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

## C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	<p>SANDEWALL E: "Towards a world-wide data base"</p> <p>COMPUTER NETWORKS AND ISDN SYSTEMS, vol. 28, no. 11, 1996, pages 1513-1522, XP004018247</p> <p>NL, NORTH HOLLAND PUBLISHING, AMSTERDAM</p> <p>ISSN: 0169-7552</p> <p>page 1513, left-hand column, line 1 - right-hand column, line 9</p> <p>page 1515, right-hand column, line 9 - page 1519, right-hand column, line 17</p>	1-17



Further documents are listed in the continuation of box C.



Patent family members are listed in annex.

\* Special categories of cited documents:

- "A" document defining the general state of the art which is not considered to be of particular relevance
- "E" earlier document but published on or after the international filing date
- "L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)
- "O" document referring to an oral disclosure, use, exhibition or other means
- "P" document published prior to the international filing date but later than the priority date claimed

- "T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention
- "X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone
- "Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art
- "&" document member of the same patent family

Date of the actual completion of the international search

13 January 2000

Date of mailing of the international search report

19/01/2000

Name and mailing address of the ISA

European Patent Office, P.B. 5818 Patentlaan 2  
NL - 2280 HV Rijswijk  
Tel. (+31-70) 340-2040. Tx. 31 651 epo nl,  
Fax: (+31-70) 340-3016

Authorized officer

Wiltink, J

# INTERNATIONAL SEARCH REPORT

International Application No

PCT/US 99/21940

## C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	<p>LINDLAND O I ET AL: "TRANSFORMATIONS IN CASE TOOLS - A COMPILER VIEW"            PROCEEDINGS OF THE INTERNATIONAL WORKSHOP ON COMPUTER AIDED SOFTWARE ENGINEERING (CASE '93), SINGAPORE,            vol. WORKSHOP 6, 19 - 23 July 1993, pages 287-296, XP000419638            LOS ALAMITOS, ca, us, IEEE COMP. SOC. PRESS            abstract            page 287, right-hand column, line 21 -page 289, right-hand column, line 5; figure 1            page 291, left-hand column, line 21 -page 292, left-hand column, line 19            page 292, right-hand column, line 4 -page 294, right-hand column, line 35</p>	1-17
A	<p>THRAMPOULIDIS K ET AL: "REDOM: AN OO LANGUAGE TO DEFINE AND ON-LINE MANIPULATE REGULATIONS IN THE RESOURCE (RE) SCHEDULING PROBLEM"            SOFTWARE PRACTICE &amp; EXPERIENCE,            vol. 27, no. 10, October 1997 (1997-10), pages 1135-1161, XP000699296            GB, JOHN WILEY &amp; SONS LTD. CHICHESTER            ISSN: 0038-0644            abstract            page 1151, line 24 -page 1154, last line; figures 7-10</p>	1-17
A	<p>DELLA MEA V ET AL: "HTML generation and semantic markup for telepathology"            COMPUTER NETWORKS AND ISDN SYSTEMS,            vol. 28, no. 11, 1996, pages 1085-1094, XP004018210            NL, NORTH HOLLAND PUBLISHING. AMSTERDAM            ISSN: 0169-7552            abstract            page 1086, right-hand column, line 43            -page 1088, left-hand column, line 25            page 1089, right-hand column, line 36            -page 1090, right-hand column, line 7</p>	1-17